

so the end result is a square-law envelope detector which removes its own second harmonics by phased cancellation rather than by fixed frequency low pass filtering.

An alternative explanation of the operational principle is that the system is a quadrature receiver (see e.g. Whalen [9]) which derives the quadrature signal from the input signal rather than from a local oscillator. A third explanation of the filter in terms of signal 'energy' and its relation to first and second derivatives of the signal has been particularly fruitful for the development of rigorous bounds on performance in relation to noise and bandwidth parameters [3].

Performance of the undersampling envelope detector: Undersampling has a simple effect: the bandwidth limits determined by Maragos et al [4] merely have to be reduced in proportion to the undersampling factor $(2n+1)$. So, for example, undersampling by a factor of 3 reduces the predicted maximum modulation bandwidth by a factor of 3. The bandwidth restrictions have been based on the signal to error ratio (SER). The error being the difference between an ideal demodulated signal and the signal demodulated by the actual filter defined by eqn. 2. Once a minimum SER has been defined then the maximum modulation bandwidth can be calculated from formulas given in [3]. The usual bandpass sampling requirements still apply [9], i.e. if B is the bandwidth then

$$B(2n+1) \leq f_c \quad (5)$$

In the schematic diagram of the quadratic filter demodulator shown in Fig. 1 the sampling frequency is simply controlled by the delay τ . To demonstrate the principle of undersampling, the outputs of such a demodulator (after performing a square root operation) have been simulated for two sampling cases. The chosen input has a Gaussian envelope and the bandwidth is about one tenth of the carrier frequency. The envelope derived by full sampling demodulation is shown in Fig. 2a. The corresponding plot for three times undersampling is shown in figure Fig. 2b. Clearly the signal to error ratio has deteriorated because the modulation bandwidth was not reduced by the requisite factor of 3. Fig. 2c shows the difference between a and b magnified by a factor of 10 for clarity. An extreme modulation bandwidth has been used to accentuate the errors. In practice, bandwidths of $<1\%$ of the carrier frequency are typical. As the SER is proportional to the square of {sampling frequency/modulation bandwidth} typical errors can be expected to be at least 20dB lower than shown.

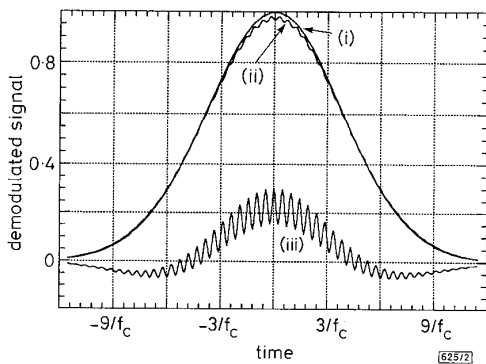


Fig. 2 Demodulated envelopes from AM signal (DSBSC) with large bandwidth Gaussian envelope

- (i) Envelope produced by nonlinear filter with full sampling, namely 4 times carrier frequency
- (ii) Envelope produced by 3 times undersampling
- (iii) $10 \times$ difference between a and b

Conclusion: A remarkably efficient nonlinear filter has been shown to work with an undersampled signal, further reducing the computational load. A direct trade-off between demodulation errors and undersampling ratio applies for signals with a given bandwidth. In situations where computational efficiency is paramount such a filter may be invaluable.

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Note on decoding binary Goppa codes

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Indexing terms: Codes, Error correction codes, Decoding

The author gives a simple expression for the polynomial $y(x)$ which solves the polynomial equation $y(x)^2 \equiv t(x) \pmod{G(x)}$, where $t(x)$, $y(x)$ and $G(x)$ are polynomials over the field $GF(2^m)$. The solution of such an equation is a step in the so called Patterson algorithm for decoding binary Goppa codes. The result may also be useful for other applications.

Introduction: For theoretical as well as practical purposes Goppa codes [3 - 5] are among the most interesting and fascinating classes of block codes. A further important application of Goppa codes stems from the introduction of the public key cryptosystem of McEliece [6]. Goppa codes can be decoded in an elegant algebraic way by using the extended Euclidean or the Berlekamp-Massey algorithm [8, 4, 5]. The Patterson algorithm for decoding binary Goppa codes ([8], algorithm 4) needs the solution of a polynomial equation of the following kind as a second step:

$$y(x)^2 \equiv t(x) \pmod{G(x)} \quad (1)$$

i.e. we are given the polynomials $t(x)$ and $G(x)$ over the field $GF(2^m)$ and want to determine the polynomial $y(x)$. In fields of characteristic two the operation of taking square roots is a linear operation, hence the usual way of taking the square root of $t(x) \pmod{G(x)}$ is to perform this linear operation using a matrix (see e.g. [8, 2] section 2.44). In the following we show that the solution of eqn. 1 can be expressed in a very simple closed form, which is very useful for direct implementation in software or hardware.

Taking square roots of polynomials modulo $G(x)$: First we set $G(x) = g_1(x)^2 + x \cdot g_2(x)^2$, hence

$$g_1(x)^2 \equiv x \cdot g_2(x)^2 \pmod{G(x)}$$

Assuming that $G(x) = G_0 + G_1x + \dots + G_nx^n$ with $G_0 \neq 0$, we know that the greatest common divisor (gcd) of x and $G(x)$ equals 1, hence there is a polynomial $w(x)$ such that $w(x)^2 \equiv x \pmod{G(x)}$.